

Compilers

Introduction

محاضرة مادة المترجمات

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للمرحلة الثالثة حاسبات /كلية التربية الاساسية

Outlines

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Language Design

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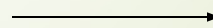
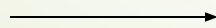
Influences on Computer Design 1.7 □

Overview and History

.Compilers are fundamental to modern computing

They act as translators, transforming human-oriented programming languages
.into computer-oriented machine languages

Programming
Language
(Source)



Machine
Language
(Target)

Overview and History (Cont'd)

The first real compiler □

FORTTRAN compilers of the late 1950s □

person-years to build 18 □

Compiler technology is more broadly applicable and has been employed in rather unexpected areas □

Text-formatting languages, like nroff and troff; preprocessor packages like eqn, tbl, pic □

Silicon compiler for the creation of VLSI circuits □

Command languages of OS □

Query languages of Database systems □

?What Do Compilers Do

Compilers may be distinguished according to the kind of target code they generate

Pure Machine Code

Augmented Machine Code

Virtual Machine Code

JVM, P-code

What Do Compilers Do? (Cont'd)

Another way that compilers differ from one another is in the format of the target machine code they generate

Assembly Language Format

Relocatable Binary Format

A linkage step is required

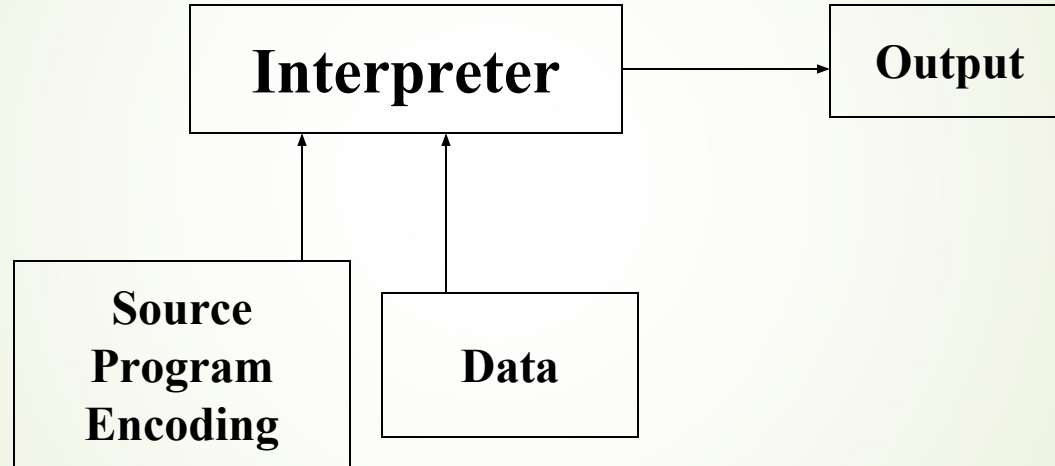
Memory-Image (Load-and-Go) Format

What Do Compilers Do? (Cont'd)

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7

Another kind of language processor, called an *interpreter*, differs from a compiler in that it executes programs without explicitly performing a translation □



Advantages and Disadvantages of an interpreter □

See page 6 & 7 □

The Structure of a Compiler

Any compiler must perform two major tasks □

Analysis of the source program □

Synthesis of a machine-language program □

The Structure of a Compiler (Cont'd)

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Source

Program

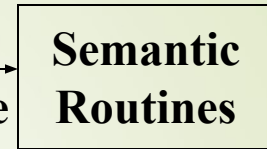
(Character Stream)



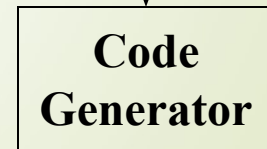
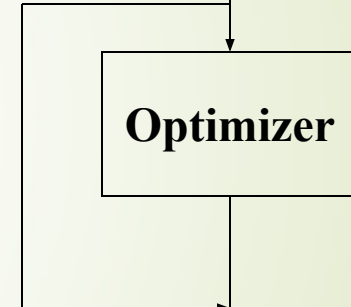
Tokens



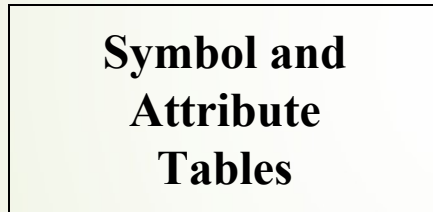
Syntactic Structure



Intermediate Representation



Target Machine Code



(Used by all Phases of The Compiler)

The structure of a Syntax-Directed Compiler

The Structure of a Compiler (Cont'd)

Scanner □

The scanner begins the analysis of the source program by reading the input, character by character, and grouping characters into individual words and symbols (**tokens**) □

The **tokens** are encoded and then are fed to the parser for syntactic analysis □

.For details, see the bottom of page 8 and page 9 □

Scanner generators □

The Structure of a Compiler (Cont'd)

Parser □

Given a formal syntax specification (typically as a context-free [CFG] grammar), the parser reads tokens and groups them into units as specified by the productions of the CFG being used □

While parsing, the parser verifies correct syntax, and if a syntax error is found, it issues a suitable diagnostic □

As syntactic structure is recognized, the parser either calls corresponding semantic routines directly or builds a syntax tree □

The Structure of a Compiler (Cont'd)

Semantic Routines □

Perform two functions □

Check the static semantics of each construct □

Do the actual translation □

The heart of a compiler □

Optimizer □

The IR code generated by the semantic routines is analyzed and transformed into functionally equivalent .but improved IR code □

This phase can be very complex and slow □

Peephole optimization □

The Structure of a Compiler (Cont'd)

One-pass compiler □

No optimization is required □

To merge code generation with semantic routines and eliminate the use of an IR □

Compiler writing tools □

Compiler generators or compiler-compilers □

E.g. scanner and parser generators □

Compiler Design and Programming Language Design

An interesting aspect is how programming language design and .compiler design influence one another □

Programming languages that are easy to compile have many advantages □

.See the 2nd paragraph of page 16 □

Compiler Design and Programming Language Design

(Cont'd)

Languages such as Snobol and APL are usually considered noncompilable

What attributes must be found in a programming language to allow compilation

Can the scope and binding of each identifier reference be determined before execution begins

Can the type of object be determined before execution begins

Can existing program text be changed or added to during execution

Compiler Classifications

- Diagnostic compilers
- Optimizing compilers
- Retargetable compiler

Compilers

Principles, Techniques, & Tools

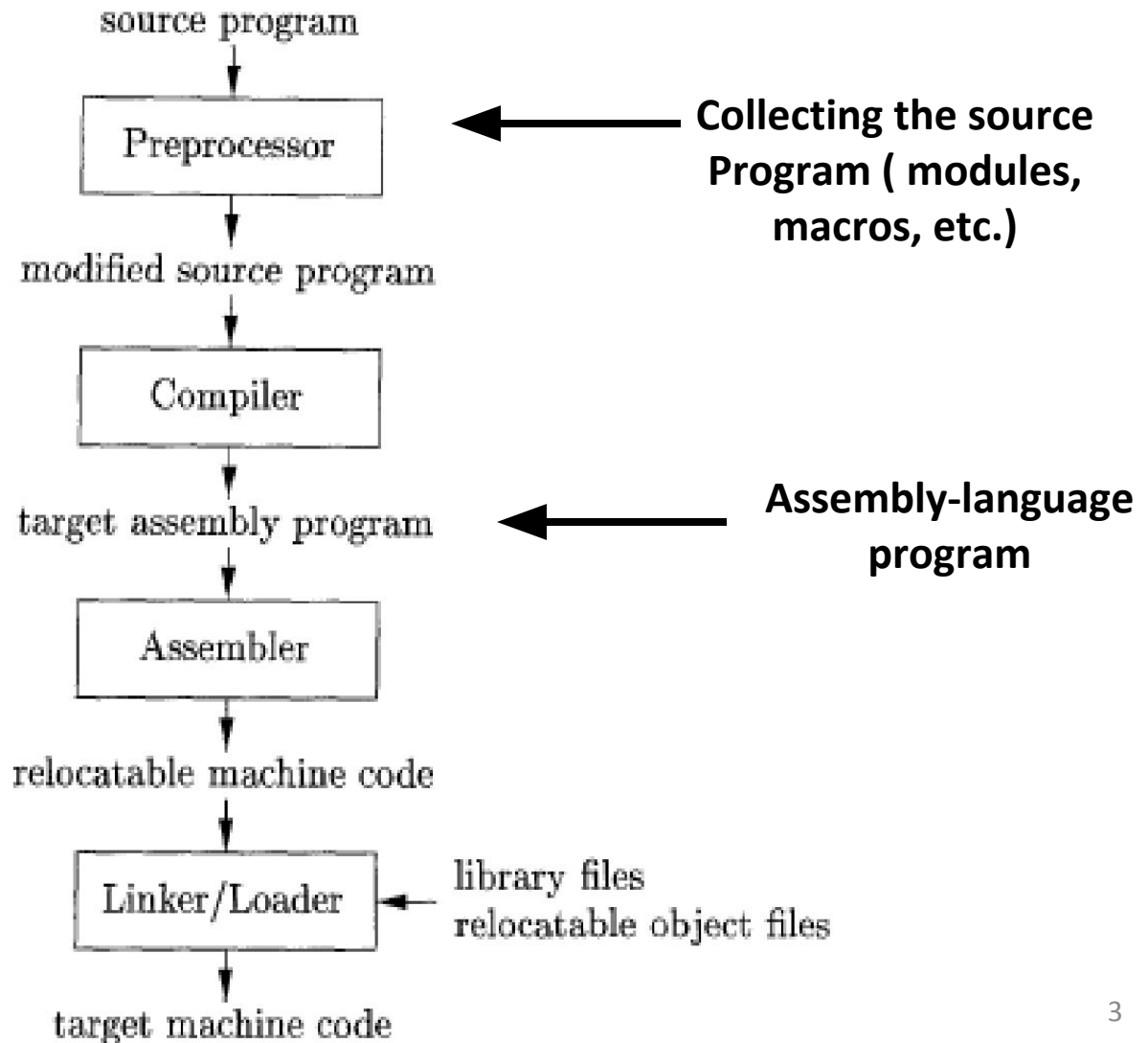
Plan

- **Introduction**
- **Lexical analysis**
- **Syntax analysis**
- **Symbol tables**

Language processing system

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To create an executable target program several programs may be required



Introduction

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- In order to reduce the complexity of designing and building computers, nearly all of these are made to execute relatively simple commands (but do so very quickly).
 - A program for a computer must be built by combining these very simple commands into a program in what is *called machine language*.
 - Most programming is done using a *high-level programming language* -> **But this language can be very different from the machine language that the computer can execute.**
- ==> This is where the compiler comes in**

Introduction

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- Using a high-level language for programming has a large impact on how fast programs can be developed:
 - Compared to machine language, the notation used by programming languages is *closer* to the way humans think about problems,
 - The compiler can spot some obvious *programming mistakes*,
 - Programs written in a high-level language tend to be *shorter* than equivalent programs written in machine language
 - The *same* program can be compiled to *many different* machine languages and, hence, be brought *to run on many* different machines.

Introduction

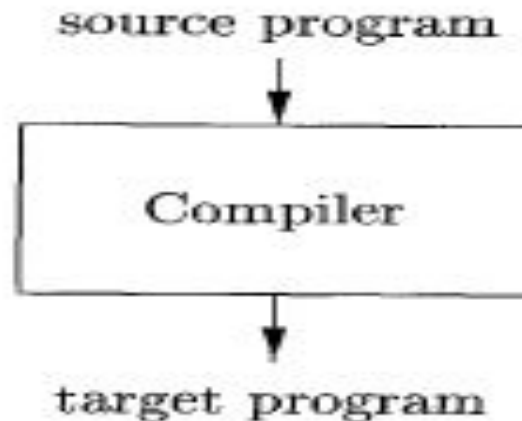
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What's a compiler ?

A compiler **translates** a program written in a high-level programming language that is suitable for human programmers into the low-level machine language that is required by computers

+

spots and **reports** obvious programmer mistakes.



Introduction

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Running the target program

- If the target program is an executable machine-language program, it can then be called by the user to process inputs and produce outputs.



Introduction

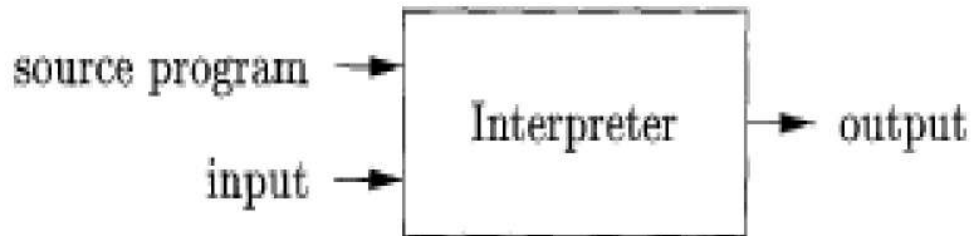
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What's an interpreter ?

- An interpreter is another way of implementing a programming language.
- Interpretation **shares** many aspects with compiling (Lexing, parsing and type-checking)

? But

Instead of producing a target program as a translation, an interpreter appears to directly execute the operations specified in the source program on inputs supplied by the user



Introduction

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Compiler vs. Interpreter

- An interpreter may need to process the same piece of the syntax tree (for example, the body of a loop) many times ☐ interpretation is **slower than** executing a **compiled program**.
- An interpreter executes the source program statement by statement ☐ it can usually give **better error diagnostics** than a compiler.

Compilers

Principles, Techniques, & Tools

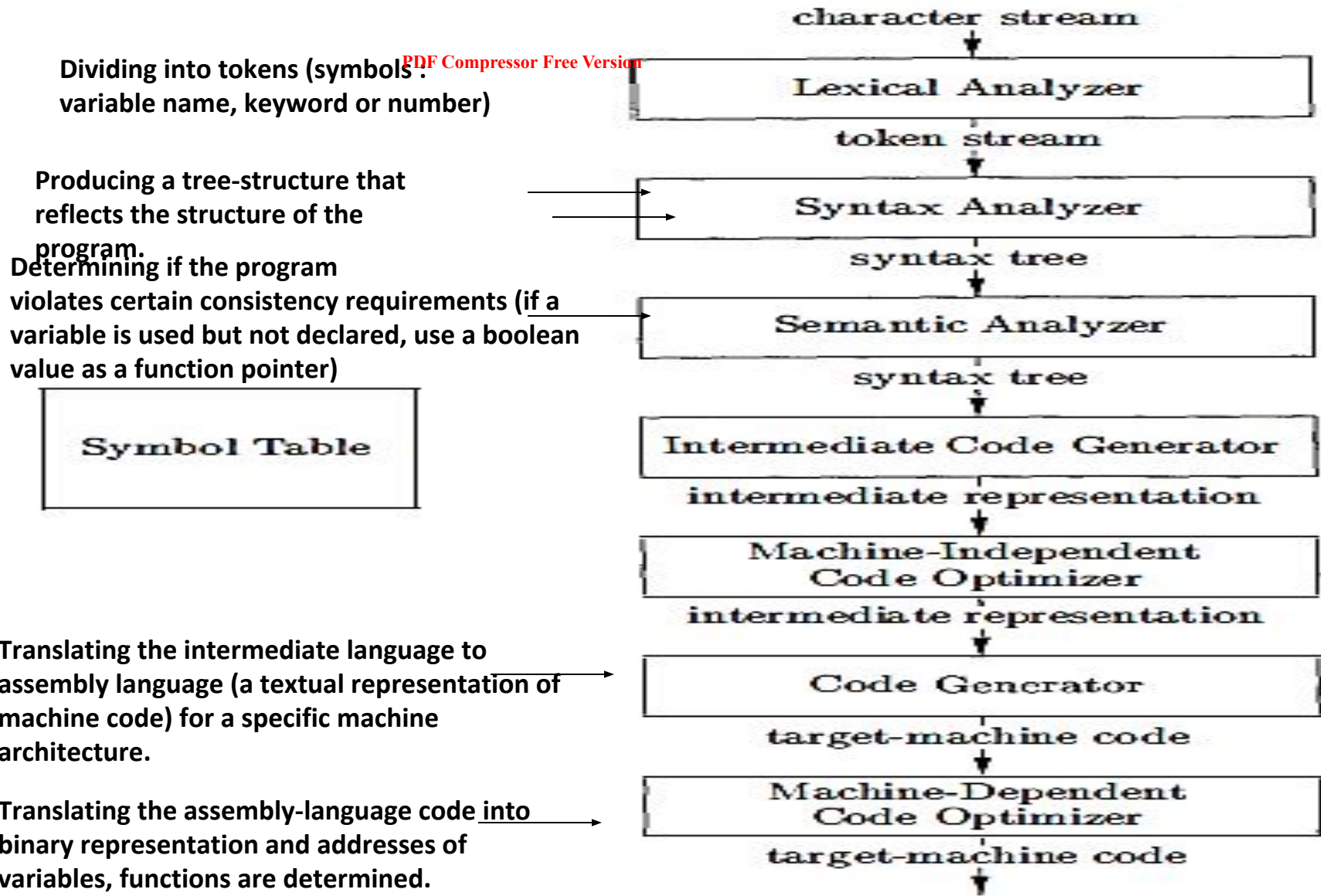
Plan

- **Introduction**
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- **Syntax analysis**
- **Symbol tables**

Introduction

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The Structure of a Compiler



Dividing into tokens (symbols: variable name, keyword or number)

Producing a tree-structure that reflects the structure of the program.

Determining if the program violates certain consistency requirements (if a variable is used but not declared, use a boolean value as a function pointer)

Symbol Table

Translating the intermediate language to assembly language (a textual representation of machine code) for a specific machine architecture.

Translating the assembly-language code into binary representation and addresses of variables, functions are determined.

Figure 1.6: Phases of a compiler

1	position	...
2	initial	...
3	rate	...

SYMBOL TABLE

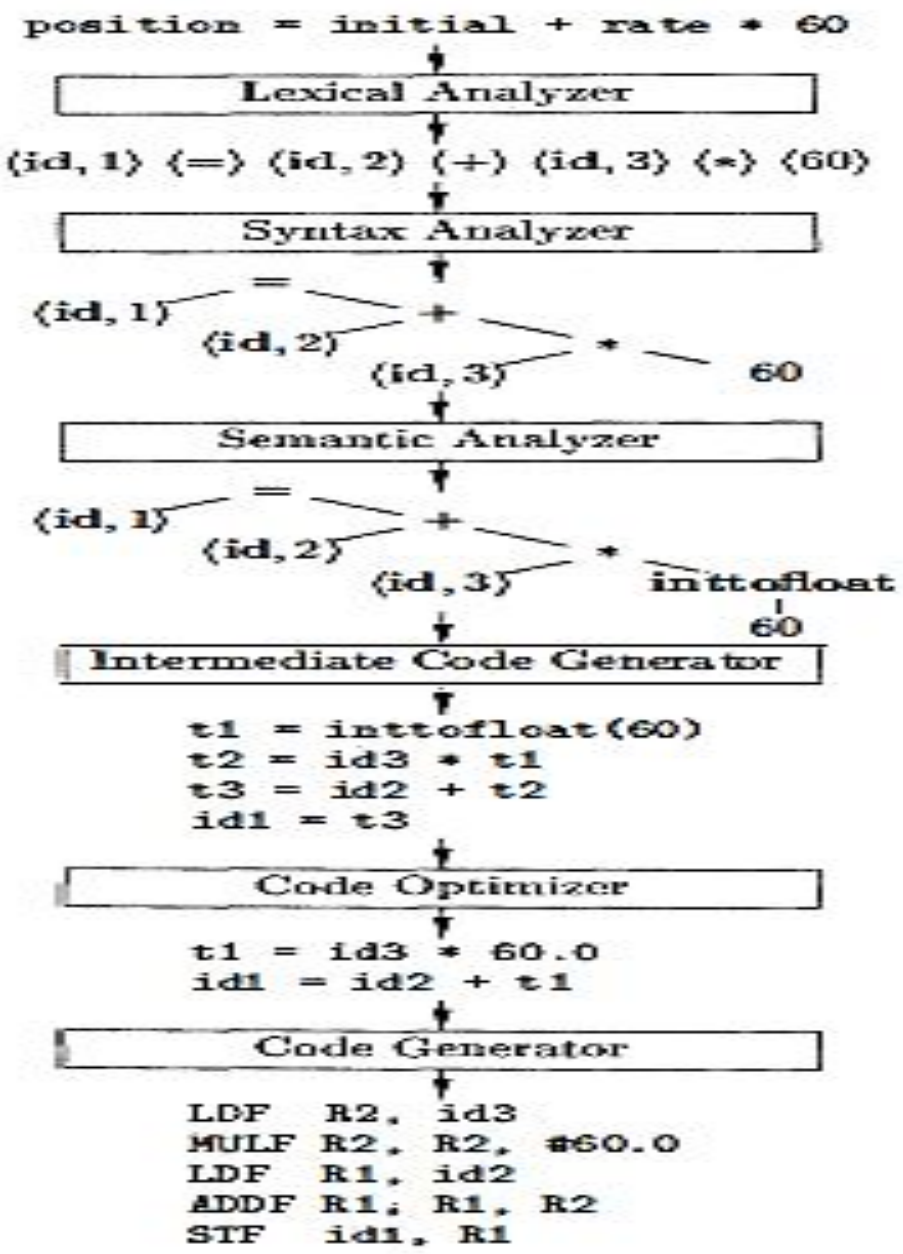


Figure 1.7: Translation of an assignment statement

Introduction

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Symbol Table Management

- The symbol table is a **data structure** containing a **record** for each **variable name**, with **fields** for the **attributes** of the name (storage allocated for a name, its type, its scope where in the program its value may be used), and for **procedure names** (number and types of its arguments, the method of passing each argument and the type returned).
- The data structure should be designed to allow **the compiler** to find the **record** for each name quickly and to **store or retrieve data** from that record quickly.

Example

Overview of the Compiler

Lexical Analysis

Break input into "*TOKENS*"

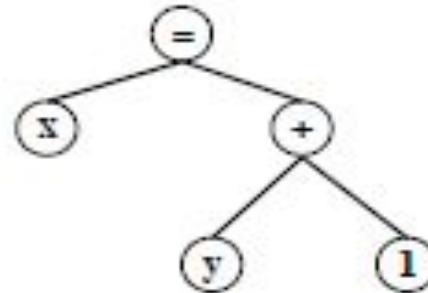
Source: `x = y + 1; /* incr x */ ...`

Tokens: `ID, EQUALS, ID, PLUS, INT, SEMI, ...`

Syntax Analysis

Context-Free Grammar

Build a parse tree



Semantic Analysis

Analyze types

Check for "semantic" errors

`x = y + true;`

Error

Introduction

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Symbol Table

One entry for each identifier

key	type	address
w	bool	50
x	int	54
y	double	58
...

Introduction

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Intermediate Code

Not machine specific

```
temp1 := x
temp2 := temp1 + 1
x := temp2
```

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Code Optimization

Eliminate redundant data movement

Optimize "goto"s to other "goto" instructions

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Code Generation

Register Assignments

Machine Specific Code

```
mov.w    x, r5
add.w    #1, r5
mov.w    r5, x
```

Introduction

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```

Error Handling

Can't just abort! ... Find more errors!

Patch things up and keep going

Lexical Errors

Syntactic Errors

Semantic Errors